

An implementation of the Volume Algorithm

Francisco Barahona and Laszlo Ladanyi

June 7, 2006

1 Introduction

Here we describe an implementation of the Volume algorithm (VA) originally presented in [1]. The following sub-directories of COIN contain the relevant pieces. The directory COIN/Vol contains the core of the algorithm. The directory COIN/Examples/VolUfl contains the necessary files for solving uncapacitated facility location problems. The directory COIN/Examples/Volume-LP contains code for dealing with combinatorial linear programs. The directory COIN/Examples/VolLP also contains code for combinatorial linear programs, this implementation relies on other parts of COIN, while the implementation in COIN/Examples/Volume-LP is self-contained. In COIN/Osi/OsiVol there is code to call the VA through OSI. The directory COIN/Examples/MaxCut contains code for doing Branch-and-Cut based on the VA, this is applied to the Max-Cut problem.

Now we give the details of each directory. We hope to receive reports about bugs and/or successful experiences.

2 COIN/Vol

This directory contains the core of the algorithm, most users should not need to modify any of the files here. The files are `INSTALL`, `Makefile` and `VolVolume.cpp`. The file `INSTALL` contains information on how to compile and build the code, on Linux it is enough to type “`make`”. The algorithm is in `VolVolume.cpp` and the header file is `VolVolume.hpp` in the directory COIN/Vol/include.

3 COIN/Examples/VolUfl

We focus here on the *uncapacitated facility location problem* (UFLP) as an example of implementation, see [3] for some of the theoretical issues. The files here are `INSTALL`, `Makefile`, `Makefile.ufl`, `ufl.cpp`, `ufl.hpp`, `ufl.par` and `data.gz`. The file `INSTALL` contains information on how to compile and build.

As a first step, a new user should be able to run the code “as is”. This can also be used as a framework for Lagrangian relaxation. The user would have to modify the files `ufl.hpp`, `ufl.cpp`, `ufl.par`, and `data`, to produce an implementation for a different problem.

Now we present the linear program used in [3]. This is

$$\min \sum c_{ij}x_{ij} + \sum f_i y_i \quad (1)$$

$$\sum_i x_{ij} = 1, \text{ for all } j, \quad (2)$$

$$x_{ij} \leq y_i, \text{ for all } i, j, \quad (3)$$

$$x_{ij} \geq 0, \text{ for all } i, j, \quad (4)$$

$$y_i \leq 1, \text{ for all } i. \quad (5)$$

Here the variables y correspond to the locations, and the variables x represent connections between customers and locations. Let u_j be a set of Lagrange multipliers for equations (2). When we dualize equations (2), we obtain the *lagrangian problem*

$$\begin{aligned} L(u) &= \min \sum \bar{c}_{ij}x_{ij} + \sum \bar{f}_i y_i + \sum u_j, \\ x_{ij} &\leq y_i, \text{ for all } i, j, \\ x_{ij} &\geq 0, \text{ for all } i, j, \\ y_i &\leq 1, \text{ for all } i. \end{aligned}$$

Where the *reduced costs* $\bar{c}_{ij} = c_{ij} - u_j$, and $\bar{f}_i = f_i$. We apply the VA to maximize $L(\cdot)$ and to produce a primal vector (\bar{x}, \bar{y}) that is an approximate solution of (1)-(5). Using this primal information we run a heuristic that gives an integer solution.

In what follows we describe the different files in this directory.

3.1 ufl.par

This file contains a set of parameters that control the algorithm and contain some information about the data. Each line has the format

keyword=value

where **keyword** should start in the first column. If we add any other character in the first column, the line is ignored or considered as a comment. The file looks as below

```
fdata=data
*dualfile=dual.txt
dual_savefile=dual.txt
int_savefile=int_sol.txt
h_iter=100

printf=3
printin=5
heurin=10

greentestin=1
yellowtestin=4
redtestin=10

lambdainit=0.1
```

```

alphainit=0.1
alphamin=0.0001
alphafactor=0.5
alphaint=50

maxsgriters=2000
primal_abs_precision=0.02
gap_abs_precision=0.
gap_rel_precision=0.01
granularity=0.

```

The first group of parameters are specific to the UFLP and the user should define them. **fdata** is the name of the file containing the data. **dualfile** is the name of a file containing an initial dual vector. If we add an extra character at the beginning (***dualfile**) this line is ignored, this means that no initial dual vector is given. **dualsavefile** is the name of a file where we save the final dual vector. If this line is missing, then the dual vector is not saved. **intsavefile** is the name of a file to save the best integer solution found by the heuristic procedure, if this line is missing, then this vector is not saved. **h_iter** is the number of times that the heuristic is run after the VA has finished.

The remaining parameters are specific to the VA. **printflag** controls the level of output, it should be an integer between 0 and 5. **printinvl=k** means that we print algorithm information every **k** iterations. **heurinvl=k** means that the primal heuristic is run every **k** iterations. **greentestinvl=k** means that after **k** consecutive green iterations the value of λ is multiplied by 2. **yellowtestinvl=k** means that after **k** consecutive yellow iterations the value of λ is multiplied by 1.1. **redtestinvl=k** means that after **k** consecutive red iterations the value of λ is multiplied by 0.67. **lambdainit** is the initial value of λ . **alphainit** is the initial value of α . **alphafactor=f** and **alphaint=k** mean that every **k** iterations we check if the objective function has increased by at least 1%, if not we multiply α by **f**.

There are three termination criteria. First **maxsgriter** is the maximum number of iterations. The second terminating criterion is as follows. **primal_abs_precision** is the maximum primal violation to consider a primal vector “near-feasible”. Let **gap_rel_precision=g**, let z be the value of the current dual solution, and p be the value of a current near-feasible primal solution. If $|z| > 0.0001$ and

$$\frac{|z - p|}{|z|} < g,$$

then the algorithms stops. Let **gap_abs_precision=f**, if $|z| \leq 0.0001$ and $|z - p| < f$ then we stop. Finally, let **granularity=k**, and let U be the value of the best heuristic integer solution found. Then if $U - z < k$ we stop.

3.2 data

The file **data** has the following format. On the first line we have the number of possible locations and the number of customers. On the next lines, the cost of opening each location appears, one cost per line. Then each of the remaining lines is like

$$i \quad j \quad d_{ij},$$

where i refers to a location, j refers to a customer, and d_{ij} is the cost of serving customer j from location i . The indices i and j start from 1. If a pair i, j is missing then the cost d_{ij} is set to 10^7 .

3.3 ufl.hpp

This file contains C++ classes specific to the UFLP.

First we have a class of parameters specific to the UFLP. The description of these parameters appears in the preceding section.

```
class UFL_parms {
public:
    string fdata;           // file with the data
    string dualfile;        // file with an initial dual solution
    string dual_savefile;   // file to save final dual solution
    string int_savefile;    // file to save primal integer solution
    int h_iter;             // number of times that the primal heuristic will be
                           // run after termination of the volume algorithm

    UFL_parms(const char* filename);
    ~UFL_parms() {}
};
```

Before the next class we should mention the classes VOL_dvector and VOL_ivector defined in VolVolume.hpp. The pseudo-code below illustrates their use.

```
int n=100;
VOL_dvector x(n); // a double vector with n entries
x=0.;             // sets to 0. all entries of x
VOL_dvector y;    // a double vector, its size remains to be set
y.allocate(n);    // size is set
y=x;              // copy each entry of x into y
VOL_dvector z(y); // a double vector of the same size as y,
                  // all entries of y are copied into z
x[0]=-1;          // first entry of x is set to -1
y[0]=x[0];        // copy first entry of x into first entry of y
```

The class VOL_ivector is used for vectors of int. One can do the same operations as for VOL_dvector.

Then we have a class containing the data for the UFLP.

```
class UFL_data { // original data for uncapacitated facility location
public:
    VOL_dvector fcost; // cost for opening facilities
    VOL_dvector dist;  // cost for connecting a customer to a facility
    VOL_dvector fix;   // vector saying if some variables should be fixed
                     // if fix=-1 nothing is fixed
    int ncust, nloc;   // number of customers, number of locations
    VOL_ivector ix;    // best integer feasible solution so far
    double      icost; // value of best integer feasible solution
public:
    UFL_data() : icost(DBL_MAX) {}
    ~UFL_data() {}
};
```

Then we have

```
class UFL_hook : public VOL_user_hooks {
public:
    // for all hooks: return value of -1 means that volume should quit
    // compute reduced costs
    int compute_rc(void * user_data,
                  const VOL_dvector& u, VOL_dvector& rc);
    // solve lagrangian problem
    int solve_subproblem(void * user_data,
                        const VOL_dvector& u, const VOL_dvector& rc,
                        double& lcost, VOL_dvector& x, VOL_dvector&v,
                        double& pcost);
    // primal heuristic
    // return DBL_MAX in heur_val if feas sol wasn't/was found
    int heuristics(void * user_data, const VOL_problem& p,
                  const VOL_dvector& x, double& heur_val);
};
```

Here the function `compute_rc` is used to compute reduced costs. In the function `solve_subproblem` we solve the lagrangian problem. In `heuristics` we run a heuristic to produce a primal integer solution.

Finally in this file we have `UFL_parms::UFL_parms(const char *filename)`, where we read the values for the members of `UFL_parms`.

3.4 ufl.cpp

This file contains several functions that we describe below.

First we have `int main(int argc, char* argv[])`. In here we initialize the classes described in `ufl.hpp`, and read the data. Then `volp.psize()` is set to the number of primal variables, and `volp.dsize()` is set to the number of dual variables. Then we check if a dual solution is provided and if so we read it.

For the UFLP all relaxed constraints are equations, so the dual variables are unrestricted. In this case we do not have to set bounds for the dual variables. If we have inequalities of the type $ax \geq b$, then we have to set the lower bounds of their dual variables equal to 0. If we had constraints of the type $ax \leq b$, then we have to set the upper bounds of their variables equal to 0. This would be done as in the pseudo-code below.

```
// first the lower bounds to -inf, upper bounds to inf
volp.dual_lb.allocate(volp.dsize);
volp.dual_lb = -1e31;
volp.dual_ub.allocate(volp.dsize);
volp.dual_ub = 1e31;
// now go through the relaxed constraints and change the lb of the ax >= b
// constrains to 0, and change the ub of the ax <= b constrains to 0.
for (i = 0; i < volp.dsize; ++i) {
    if ("constraint i is '<=' ") {
        volp.dual_ub[i] = 0;
    }
}
```

```

    if ("constraint i is '>=' ") {
        volp.dual_lb[i] = 0;
    }
}

```

The function `volp.solve` invokes the VA. After completion we compute the violation of the fractional primal solution obtained. This vector is `psol`. Then we check if the user provided the name of a file to save the dual solution. If so, we save it. Then we run the primal heuristic using `psol` as an input. Notice that this heuristic has also been run periodically during the execution of the VA. Then if the user has provided the name of a file to save the integer heuristic solution, we do it. Finally the values of the solutions and some statistics are printed.

The next function is `void UFL_read_data(const char* fname, UFL_data& data)`, where we read the data. `data.nloc` is the number of locations, `data.ncust` is the number of customers. `data.fcost` is a vector containing the cost of opening each location. `data.dist` is a vector containing the cost of serving customers from facilities. All entries are initialized to 10^7 and then particular entries are being set with the statement

```
dist[(i-1)*ncust + j-1]=cost;
```

where `i` is the index of a location and `j` is the index of a customer. Here the indices start from 1. Finally we have a vector `data.fix` associated with the locations. A particular entry is set to 0 if the location should be closed, it is set to 1 if it should be open, and it is set to -1 if this variable is free. Initially all entries are set to -1.

In the function

```
double solve_it(void * user_data, const double* rdist, VOL_ivector& sol)
```

we solve the lagrangian problem. We receive the data and reduced costs as input and return a primal vector. The solution is in the vector `sol`. Its first `n` entries correspond to the locations, then all remaining entries correspond to connections between locations and customers.

In the function

```
int UFL_hook::compute_rc(void * user_data, const VOL_dvector& u, VOL_dvector& rc)
```

we compute the reduced costs. They will be used to solve the lagrangian problem.

In the function

```
int
UFL_hook::solve_subproblem(void *user_data,
                           const VOL_dvector& u, const VOL_dvector& rc,
                           double& lcost, VOL_dvector& x,
                           VOL_dvector& v, double& pcost)
```

we compute the lagrangian value, we call `solve_it`, we compute the objective value and the vector `v` defined as follows. If \hat{x} is the primal solution given by `solve_it`, and $Ax \sim b$ is the set of relaxed constraints, then the difference `v` is

$$v = b - A\hat{x}.$$

The last function in this file is

```
int
UFL_hook::heuristics(void * user_data, const VOL_problem& p,
                     const VOL_dvector& x, double& icost)
```

where we run the following simple heuristic. Given a fractional solution (\bar{x}, \bar{y}) , let \bar{y}_i be the variable associated with location i . We pick a random number $r \in [0, 1]$ and if $r < \bar{y}_i$ facility i is open, and

closed otherwise. We repeat this for every facility, then given the set of open facilities we find a minimum cost assignment of customers. This function is invoked periodically in the VA and by the main program after the VA has finished.

4 COIN/Examples/Volume-LP

Here we focus on *Combinatorial Linear Programs*, these are linear programs where the matrix has 0, 1, -1 coefficients and the variables are bounded between 0 and 1. The VA has been very effective at producing fast approximate solutions to these LPs, see [2]. As a first step, a new user should be able to run our code “as is”. The input should be an MPS file.

Initially this directory contains the files: `README`, `Makefile`, `lp.hpp`, `lp.cpp`, `lp.par`, `data.mps.gz`, `lpc.h`, `lpc.cpp`, `reader.h`, `reader.cpp`. On a Unix system one should type “make”, “gunzip data.mps.gz” and “volume-lp” to run the code. Then the code will run and produce the files `primal.txt` and `dual.txt` that contain approximate solutions to both the primal and the dual problem.

We assume that we have an LP like

$$\min cx \tag{6}$$

$$Ax \sim b \tag{7}$$

$$l \leq x \leq u. \tag{8}$$

Let π be a set of Lagrange multipliers for constraints (7). When we dualize them we obtain the *lagrangian problem*

$$\begin{aligned} L(u) &= \min (c - \pi A)x + \pi b, \\ l &\leq x \leq u. \end{aligned}$$

We apply the VA to maximize $L(\cdot)$ and to produce a dual vector $\bar{\pi}$, and primal vector \bar{x} that is an approximate solution of (6)-(8).

In what follows we describe the files `lp.par` and `data.mps`.

4.1 lp.par

This file contains a set of parameters that control the algorithm and contain information about the data. Each line has the format

keyword=value

where **keyword** should start in the first column. If we add any other character in the first column, the line is ignored or considered as a comment. The file looks as below

```
fdata=data.mps
*dualfile=dual.txt
dual_savefile=dual.txt
primal_savefile=primal.txt
h_iter=0
var_ub=1.0
```

```

printflag=3
printinvl=20
heurinvl=100000000

greentestinvl=2
yellowtestinvl=2
redtestinvl=10

lambdainit=0.1
alphainit=0.01
alphamin=0.0001
alphafactor=0.5
alphaint=80

maxsgriters=2000
primal_abs_precision=0.02
gap_abs_precision=0.
gap_rel_precision=0.01
granularity=0.

```

The first group of parameters are specific to LP and the user should define them. `fdata` is the name of the file containing the data. `dualfile` is the name of a file containing an initial dual vector. If we add an extra character at the beginning (`*dualfile`) this line is ignored, this means that no initial dual vector is given. `dual_savefile` is the name of a file where we save the final dual vector. If this line is missing, then the dual vector is not saved. `primal_savefile` is the name of a file to save the primal solution, if this line is missing, then this vector is not saved. `h_iter` is the number of times that the heuristic is run after the VA has finished. We did not include a heuristic in this implementation. `var_ub` is an upper bound for all primal variables, for 0-1 problems we set `var_ub=1`.

The remaining parameters are specific to the VA. `printflag` controls the level of output, it should be an integer between 0 and 5. `printinvl=k` means that we print algorithm information every k iterations. `heurinvl=k` means that the primal heuristic is run every k iterations. `greentestinvl=k` means that after k consecutive green iterations the value of λ is multiplied by 2. `yellowtestinvl=k` means that after k consecutive yellow iterations the value of λ is multiplied by 1.1. `redtestinvl=k` means that after k consecutive red iterations the value of λ is multiplied by 0.67. `lambdainit` is the initial value of λ . `alphainit` is the initial value of α . `alphafactor=f` and `alphaint=k` mean that every k iterations we check if the objective function has increased by at least 1%, if not we multiply α by f .

There are three termination criteria. First `maxsgriter` is the maximum number of iterations. The second terminating criterion is as follows. `primal_abs_precision` is the maximum primal violation to consider a primal vector “near-feasible”. Let `gap_rel_precision=g`, let z be the value of the current dual solution, and p be the value of a current near-feasible primal solution. If $|z| > 0.0001$ and

$$\frac{|z - p|}{|z|} < g,$$

then the algorithms stops. Let `gap_abs_precision=f`, if $|z| \leq 0.0001$ and $|z - p| < f$ then we stop.

Finally, let `granularity=k`, and let U be the value of the best heuristic integer solution found. Then if $U - z < k$ we stop. We did not include any heuristic in this implementation.

4.2 data.mps

This is an MPS file that is read with code in `reader.cpp`. If a different type of input has to be used, one should change the code in `reader.cpp`.

5 COIN/Examples/VolLp

This code treats linear programs in a similar way as it is done in `COIN/Examples/Volume-LP`, the main difference is that this code uses other components of `COIN-OR` while the code in

`COIN/Examples/Volume-LP` is self contained. The files here are `INSTALL`, `Makefile`, `Makefile.vollp`, `vollp.cpp`. The `INSTALL` contains instructions on how to compile and build the code. The input should be an MPS file.

6 COIN/Osi/OsiVol

In this directory there is the file `OsiVolSolverInterface.cpp` that allows the user to call the VA through `OSI`. This is also intended to deal with combinatorial linear programs. The code below reads an MPS file and calls the VA through `OSI`.

```
#include "OsiVolSolverInterface.hpp"

#include <iostream>

int main(int argc, char *argv[])
{
    OsiVolSolverInterface osilp;

    osilp.readMps("file","mps");

    osilp.initialSolve();

    const int numCols=osilp.getNumCols();
    const double *x=osilp.getColSolution();
    for (int j=0; j<numCols; ++j){
        std::cout << j << " " << x[j] << "\n";
    }

    return 0;
}
```

7 COIN/Examples/MaxCut

This directory contains code for doing Branch-and-Cut based on the VA as in [4]. The code is specialized to the Max-Cut problem.

References

- [1] F. BARAHONA AND R. ANBIL, *The volume algorithm: producing primal solutions with a sub-gradient method*, Math. Program., 87 (2000), pp. 385–399.
- [2] ———, *On some difficult linear programs coming from set partitioning*, Discrete Appl. Math., 118 (2002), pp. 3–11. Third ALIO-EURO Meeting on Applied Combinatorial Optimization (Erice, 1999).
- [3] F. BARAHONA AND F. A. CHUDAK, *Solving large scale uncapacitated facility location problems*, in Approximation and complexity in numerical optimization (Gainesville, FL, 1999), vol. 42 of Nonconvex Optim. Appl., Kluwer Acad. Publ., Dordrecht, 2000, pp. 48–62.
- [4] F. BARAHONA AND L. LADANYI, *Branch-and-cut based on the volume algorithm: Steiner trees in graphs and max-cut*, Report, 2001. available at <http://optimization-online.org>.